

This exam consists of 6 exercises, from which 100 points can be obtained in total. The division of the points over the various parts is as follows:

Question:	1	2	3	4	5	6	Total
Points:	25	15	20	20	10	10	100
Score:							

Exam grade: $\frac{\text{total number of points}}{10}$

Final course grade: $\frac{1}{4}$ Weekly pretest grade + $\frac{3}{4}$ Exam grade

1. Consider the following LP, which we will refer to as the “primal LP”.

$$\begin{aligned}
 \max \quad & z = 4x_1 - x_2 + 3x_3 \\
 \text{s.t.} \quad & 2x_1 + 2x_2 + x_3 + x_4 = 5 \\
 & x_1 + 3x_2 + x_3 - x_4 \leq 15 \\
 & x_1, x_2, x_3, x_4 \geq 0
 \end{aligned}$$

- (a) (10 points) Determine the dual of this LP.
- (b) (5 points) The primal LP has been solved using the simplex method, yielding the final tableau below.

Basic	z	x_1	x_2	x_3	x_4	s_2	value
z	1	2	7	0	3	0	15
x_3	0	2	2	1	1	0	5
s_2	0	-1	1	0	-2	1	10

Write down the primal optimal solution and objective value. Also state the objective value of an optimal dual solution (you don't need to give the actual optimal solution to the dual). Explain why there is a difference in the variables listed in in the tableau compared to the original LP.

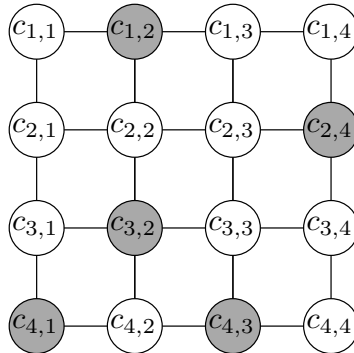
- (c) (10 points) Now introduce the new constraint

$$x_1 + 2x_3 \leq 6,$$

in addition to the existing constraints in the primal LP. Determine an optimal solution to the new LP (along with the corresponding new objective value) bu using the dual simplex method.

(Hint: if you do the calculations correctly, all fractions you see will be multiples of $\frac{1}{3}$, and only one pivot step will be required.)

2. (15 points) You want to build some windmills for generating power out in the North Sea. The possible windmill locations form a grid, as shown. Different locations will incur different costs, because of variations in the sea floor; we denote the cost of site at row i and column j by $c_{i,j}$, as shown.



You must build exactly 5 windmills. However, it is not allowed to place two windmills on two sites that are directly adjacent, either horizontally or vertically, because this will cause wind interference. The grayed circles show one possible solution that is valid. (Notice that you can place two windmills on squares that are diagonal from each other.)

Formulate the problem of building all the required windmills as an integer linear program. Explain the meaning of all your variables and constraints.

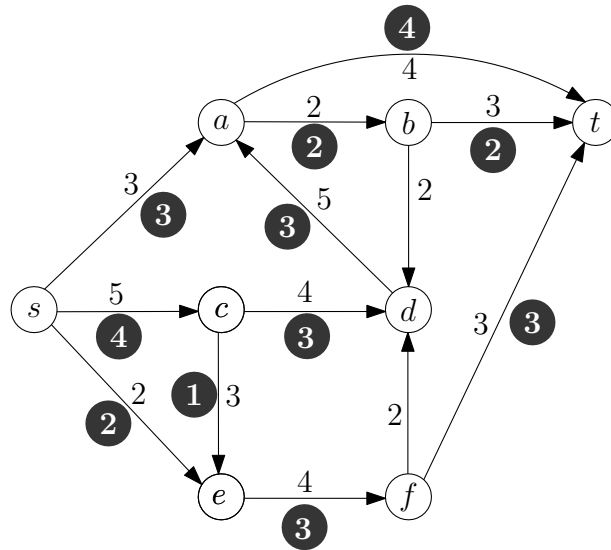
3. Consider the following problem. You are packing your suitcase for an overseas trip, and need to decide what items to bring. You have 4 items (just one of each); each has a certain value to you, and a certain weight:

$$\begin{array}{r|cccc} \text{weight (kg)} & 3 & 5 & 4 & 1 \\ \text{value} & 7 & 12 & 10 & 2 \end{array}$$

You can only carry at most 10kg of weight. The problem is to decide which items you should put into your knapsack to maximize the overall value.

- (5 points) Write down an ILP for this problem.
- (5 points) Determine the unique optimal solution of the LP relaxation of this problem. Also give the corresponding optimal objective value.
- (10 points) Solve the ILP from (a) using Branch and Bound. Clearly indicate in which order you compute the nodes of the search tree and where you prune the search tree, and based on which pruning criterion.

4. Consider the following directed graph; arc capacities are shown. In addition, an s - t -flow f has been given (the circled numbers).



- (a) (2 points) Write down the value of the flow f .
- (b) (8 points) Determine either i) a flow of strictly larger value, or ii) argue convincingly that f is a maximum flow.
- (c) (10 points) Find a shortest s - t -path in the digraph using Dijkstra's algorithm, assuming that every arc has length 1 (do *not* use the capacities of the arcs as lengths!). Show your working.
5. (a) (4 points) You are considering two algorithms for solving a certain problem. On an input of size n , the worst-case running time of the first algorithm is $n^2 + n2^n + 100$, and of the second algorithm is $100n^4 + 200n^3 + 3000$. Which algorithm is faster, assuming a large input size? Explain your answer.
- (b) (6 points) What is the length of a longest increasing subsequence of the following sequence?

(1, 3, 2, 6, 4, 10, 7, 9).

You *must* solve this problem using dynamic programming – show your working. (You do not need to provide the actual subsequence, but it may be a good idea for checking your answer)

6. (10 points) **This question is challenging; I recommend completing the rest of the exam before attempting it.**

Consider the following variation of the edit distance, which I'll call the modified edit distance. One can add a letter, or modify a letter, as before. But now, one can delete any sequence of consecutive letters, and this counts as just one operation. For example: given the initial word REASON and final word WRONG, one could delete EAS to get RON, then add W to get WRON, and then add G to get WRONG. This is in fact the best solution, so the modified edit distance is 3.

Thinking of this in terms of columns: the configuration is

$$\begin{array}{c|c|c|c|c|c|c|c|c} - & R & E & A & S & O & N & - \\ W & R & - & - & - & O & N & G \\ 1 & & & 1 & & & & 1 \end{array}$$

But instead of paying 1 for each column with a change, we pay only 1 for columns 3-5.

Describe a dynamic programming algorithm to determine the modified edit distance, given an initial word $x_1x_2\dots x_m$ that must be changed to $y_1y_2\dots y_n$. You only need to determine the modified edit distance itself, not the way that the word should be transformed.

You should use the same subproblems as with the original edit-distance problem; thus, define

$$E(i, j) = \text{modified edit distance between } x_1x_2\dots x_i \text{ and } y_1y_2\dots y_j.$$

In fact, you can use the following template. All you need to do is decide what should go in the placeholders [[1]], [[2]] and [[3]] (with explanation).

- 1: $E(0, 0) = 0$
- 2: $E(i, 0) = [[1]]$ for all $1 \leq i \leq m$
- 3: $E(0, j) = [[2]]$ for all $1 \leq j \leq n$
- 4: **for** $i = 1, 2, \dots, m$ **do**
- 5: **for** $j = 1, 2, \dots, n$ **do**
- 6: $E(i, j) = [[3]]$
- 7: **return** $E(m, n)$